

Catch Me (If You Can): Data Survival in Unattended Sensor Networks

New threats and countermeasures

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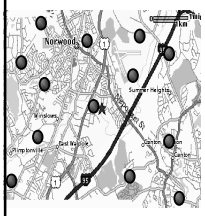
PerCom 2008, Hong Kong

Agenda

- Introduction
  - A different kind of WSN
  - New adversarial model
- Naïve defense strategies
- Cryptography to the rescue
- Related Work
- Conclusions and Future directions

A “Typical” Wireless Sensor Network







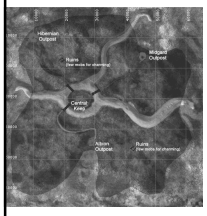
Many real, alleged and imagined applications

- Networking
  - Sensor-to-sink communication (opt. sink-to-sensors)
- Collection method
  - Periodic collection
    or
  - Event driven
    or
  - Query based = on-demand
- Online Sink
  - Real-time off-loading of data

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Unattended
Wireless Sensor Network (UWSN)





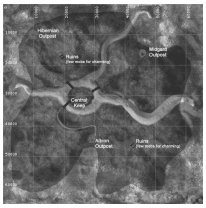


- Nodes operate in hostile environment
  - Initial deployment might be ad-hoc
- No ever-present sink
  - Itinerant
- Periodic data sensing
  - Nodes might retain data for a long time
  - Data might be valuable
- Nodes are left on their own
  - Adversary roams around with impunity
  - **Challenge: Data Survival in UWSNs**

Examples



- WSN deployed in a recalcitrant country to monitor any potential nuclear activity



- Underground WSN monitoring sound and vibration produced by troop movements or border crossings



- Anti-poaching WSN in a national park tracking/recording firearm discharge locations

New kind of Adversary (Adv)

- Mobile
  - Migrates between sets of nodes between sink visits
- Well-informed
  - Knows network topology and network defense strategy
- Erratic (seemingly)
  - Unpredictable and possibly untraceable movements
- Data-centric
  - No interference with sensing or network operation
- Powerful (but not omnipotent)
  - Compromises up to a certain # of nodes

Assumptions

- Scheduled (per round) data sensing/collection
  - Max ***v*** rounds between sink visits
- Adv compromises at most ***k*** (out of ***n***) nodes per round
  - Compromised nodes not necessarily contiguous
  - Reads all storage
  - Listens to all incoming and outgoing communication
- Adv knows which data to target and when it was sensed
  - Receives external signal at collection time (round 0)
    - Target node identity + collection round
    - Possibly, also target value
- UWSN knows nothing…
  - Equal protection for all data

What if sensors have no crypto capability?

- Cheap sensors
  - No crypto
  - Can only (attempt to) hide data location
- Data Migration strategies
  - Do Nothing
  - Move Once
  - Keep Moving
- Adv Goal: Search-and-erase
  - Looks for target data in compromised sensors
- Adv strategy:
  - Lazy
  - Frantic
  - Smart

Survival vs. Attack Strategies

	Attack Strategy		
Survival Strategy	LAZY	FRANTIC	SMART
DO NOTHING	✗	✓	✗
MOVE ONCE	✗	✓	✗
KEEP MOVING	✓	✓	✓

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Do Nothing

- Data kept at originating sensor
  - Trivial
- Adversary wins in one round
  - Round 0
    - Learns originating sensor
  - Round 1
    - Corrupts it
    - Deletes target data

Move Once

- Data off-loaded to a random recipient node
  - Kept there for all subsequent rounds
- Adversary wins in at most $\left\lceil \frac{n}{k} \right\rceil$ rounds
  - Round 0
    - Learns originating node
    - Data not there anymore
  - Round i
    - Move to next set of previously uncompromised nodes
  - At most $\left\lceil \frac{n}{k} \right\rceil$ rounds to find and erase it

Keep Moving

Algorithm 1: KEEP-MOVING

```
   /* start round 0 */
   all nodes sense their values
   each node exchanges data with others
0  A learns s̄_0 and x                          ← Adv learns target data at round 0
   /* end round 0 */
   SET z=min(v, n/k)
   SET found=FALSE
   for ((r ← 1 to z) and (not found)) do
      /* start round r */
1     select C_r  /* new set of nodes to compromise */
2     compromise C_r and release C_{r-1}
3     if (x found on some s_i ∈ C_r) then      ← Adv looks for target data in the new set of compromised nodes
3.1      delete x
3.2      SET found=TRUE
      all nodes sense their values
      each node exchanges data with others     ← Nodes exchange messages
4     if (x received by some s_i ∈ C_r) then   ← Adv looks for target data in the messages received by corrupted nodes
4.1      delete x
4.2      SET found=TRUE
      /* end round r */
```

- Adv can have up to two chances per round
  - Before data exchange
  - After data exchange

Keep Moving – Lazy

- Exploit the fact that data is constantly moving among sensors
- Two chances at round 1; one chance each new round

- Prob. data survives v rounds

$$P_L(v) = P_1 \cdot P_2^{v-1}$$

$$P_1 = \frac{k}{n} + \left(1 - \frac{k}{n}\right)\frac{k}{n} = \left(1 - \frac{k}{n}\right)^2$$

$$P_2 = 1 - \frac{k}{n}$$

Keep Moving – Frantic

- Select a new random set of sensors to compromise at each round
- Two chances per round
- Prob. data survives v rounds

$$P_F(v) = P_1 \cdot P_2^{v-1} \cdot P_3^{v-1}$$

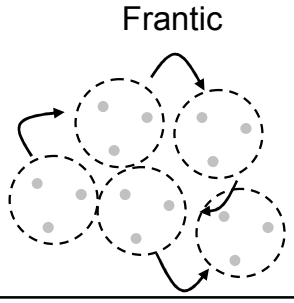
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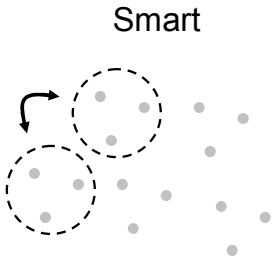
$$P_2 = 1 - \frac{k}{n}$$

$$P_3 = 1 - \frac{k}{n-k}$$

Keep Moving – Smart

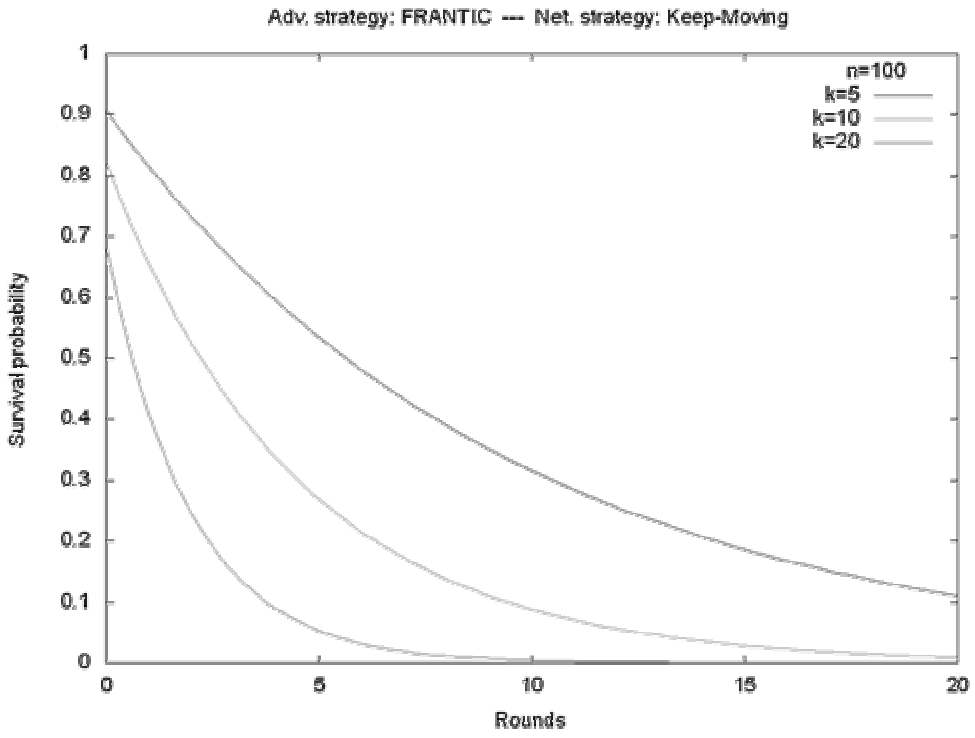
- Moves between two fixed (non-overlapping) set of nodes
  - No matter adversarial strategy, data recipient node is always chosen according to an uniform distribution
  - Same success probability!



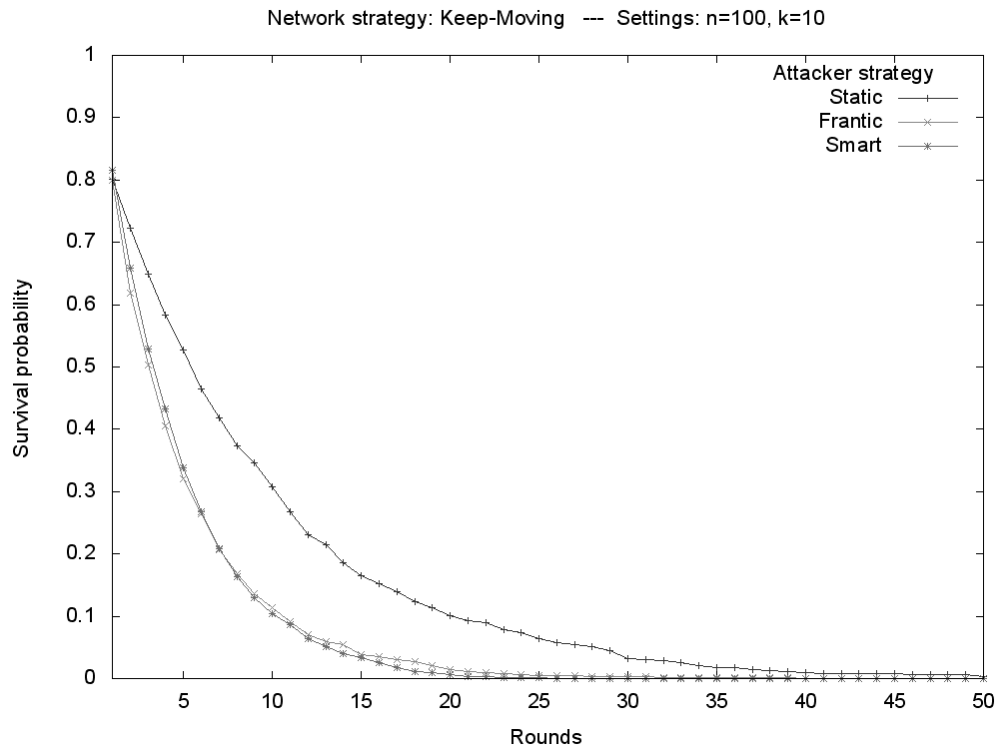


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Results


Adv. strategy: FRANTIC --- Net. strategy: Keep-Moving
n=100
k=5
k=10
k=20
Survival probability
Rounds
0
0.1
0.2
0.3
0.4
0.5
0.6
0.7
0.8
0.9
1
0
5
10
15
20

Keep Moving – Smart



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Overhead 1

scheme	*msg × round (c)*	*msg tot*	*Queue*	*msg rec*
Move-Once	n	$v \cdot n$	$Pr[\exists\, s_i \text{ s.t. } L_i^r \geq r + \sqrt{nr}] \leq e^{-r/2+\ln n}$	$O(\ln n)$
Keep-Moving	$r \cdot n$	$(r^2/2)n$	$Pr[\exists\, s_i \text{ s.t. } L_i^r \geq 2er] \leq 2^{-r+\ln n}$	$Pr[\exists\, s_i \text{ s.t. } M_i^r \geq 2er] \leq 2^{-r+\ln n}$

- Prob. # stored messages do not exceeds a given value ℓ
- L^i_r = # msg stored on s_i at round r $\rightarrow$ $E[L_i^r] = r$
- From the method of bounded differences, given $\ell > r + \sqrt{rn}$

$$Pr[L_1^r \geq \ell \cup \ldots \cup L_n^r \geq \ell] \leq nPr[L_1^r \geq \ell] \leq e^{-r/2+\ln n}$$

Overhead 2

scheme	*msg × round (r)*	*msg tot*	*Queue*	*msg rec*
Move-Once	n	$\sim n$	~~$Pr[\exists\, s_i \text{ s.t. } L_i^r \geq r+\sqrt{nr}] \leq e^{-r/2+\ln n}$~~	$O(\ln n)$
Keep-Moving	$r \cdot n$	$(r^2/2)n$	$Pr[\exists\, s_i \text{ s.t. } L_i^r \geq 2er] \leq 2^{-r+\ln n}$	$Pr[\exists\, s_i \text{ s.t. } M_i^r \geq 2er] \leq 2^{-r+\ln n}$

- Prob. # stored messages do not exceeds a given value ℓ
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$$Pr[L_1^r \geq \ell \cup \ldots \cup L_n^r \geq \ell] \leq nPr[L_1^r \geq \ell] \leq e^{-r/2+\ln n}$$

- Variables L^i_r are independent → Chernoff bound → $Pr[L_1^r \geq \ell] \leq 2^{-r}$ for $\ell > 2er$
- M^i_r = # msg received by s_i at round r

$$Pr[M_1^r \geq \ell \cup \ldots M_n^r \geq \ell] \leq nPr[M_1^r \geq \ell] \leq 2^{-r+\log_2 n}$$

Replication

- Each sensor produces R copies of its reading
  - Information survives as long as one copy survives
- $X_{i,j} = 1$ if $Pr[X_{1,j} = 1] = P_1 \cdot P_2^{j-1} \cdot P_3^{j-1}$ to round j

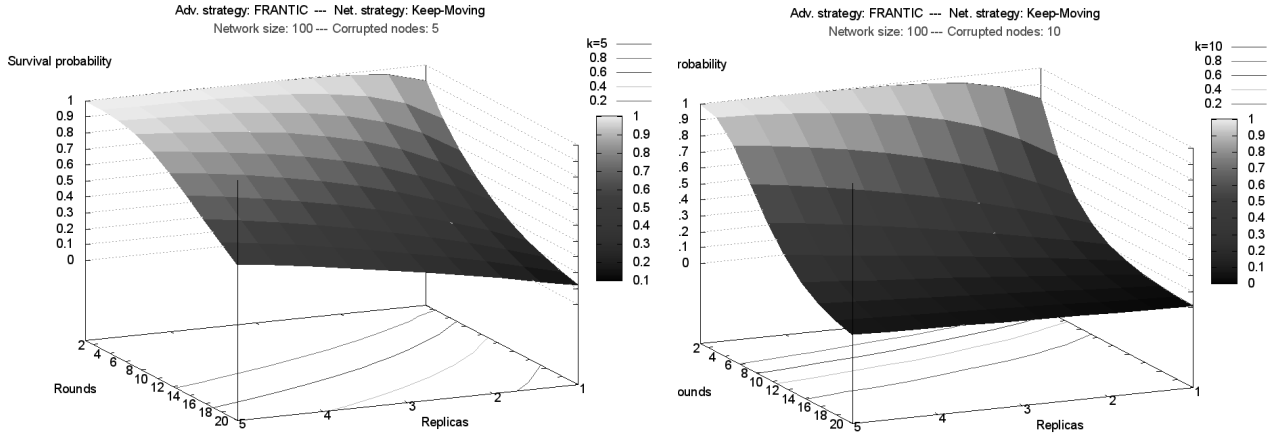
$$\overline{P_R^v} = Pr[X_{1,v} = 0 \wedge \ldots \wedge X_{R,v} = 0] = Pr[X_{1,v} = 0]^R =$$
$$= (1 - Pr[X_{1,v} = 1])^R = (1 - P_1 \cdot P_2^{v-1} \cdot P_3^{v-1})^R$$

- Prob. information survives:

$$P_R^v = 1 - \overline{P_R^v} = 1 - (1 - P_1 \cdot P_2^{v-1} \cdot P_3^{v-1})^R$$

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Results



- Replication of sensed data
  - Increases survival probability
  - Requires more storage and power
  - Given enough rounds, **Adv always wins!**

Encryption

- Goal: hide data contents and origin from the adversary
- Adv can not decrypt



- Adv can not identify data to erase

- Randomized Encryption
  - Random values involved in the encryption process
  - Given two ciphertexts encrypted under the same key, it is infeasible to determine whether the corresponding plaintexts are the same

Public Key Encryption

- Each node knows sink's public key PK_S
- d_i^r -- data sensed by s_i at round r stored as

$$E_i^r = E(PK_S, r, s_i, etc.)$$

- Adv can only try to brute-force guessing the plaintext
  - If random data is involved in the encryption, ciphertext guessing becomes infeasible (i.e., randomized encryption)

Symmetric Key Encryption

- s_i shares k_i^0 with the sink
- d_i^r -- data sensed by s_i at round r is stored as

$$E_r^i = E(k_r^i, d_r^i, r, etc.)$$

- Forward security
  - per round key evolution

$$k_{r+1}^i = OWF(k_r^i)$$

- Adv can not compute previous keys

Related Work

- Mobile Ad Hoc Networks
  - Data availability in partitioned MANETs
    - [Hara , et al. 2006, Giannuzzi, et al. 2005]
  - Multi-path routing to improve confidentiality and availability
    - [Papadimitratos, et al. 2006, Berman, et al. 2005]
- Sensor Networks
  - Data coding to increase data recovery in presence of disasters
    - [Kamra, et al. 2006]

Conclusion and Future Directions

- Contribution:
  - New kind of network - UWSN
  - New mobile UWSN adversary
  - Simple solutions to provide data survival
- Future work:
  - Explore the design space of crypto
    - Encryption
    - Authentication
  - New adversarial models
    - What if the adversary is *proactive?*
    - Different goals

Questions?

Thank you!